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## Availability analysis of two hosts systems

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### Abstract

This paper analysis two host system prevalent in banking industries using RPGT for system parameters considering constant mean time to failure and repair for hardware and software followed by illustration and special cases.

**Keywords:** Availability, base-state, regenerative point graphical technique (RPGT), system parameters

### Introduction

This paper analyses two host system prevalent in banking industry is using RPGT for system parameters MTSF, Availability, Busy Period of the Server Visits, Expected Number of Server Visits considering constant mean time to failure and repair for hardware and software. Considering the importance of individual units in system Kumar, J. & Malik, S. C. [1] have discussed the concept of preventive maintenance for a single unit system. Liu. R. [2], Malik, S.C. [3], Nakagawa, T. and Osaki, S. [4] have discussed reliability analysis of a one unit system with un-repairable spare units and its applications. Goel, P. & Singh J. [5], Gupta, P., Singh, J. & Singh, I.P. [6], Kumar, S. & Goel, P. [7], Gupta, V. K. [8], Chaudhary, Goel & Kumar [9] Sharma & Goel [10], Ritikesh & Goel [11] and Goyal & Goel [12] have discussed behavior of systems with perfect and imperfect switch-over of systems using various techniques.

**Model Description:** Two hosts system is the availability analysis of a banking industry having two hosts is considered. In a bank branch where there are two hosts, each host have a software (S) operated by a hardware (S). Here, in the present study there are two hosts with software  $S_1, S_2$  operated by hardware  $H_1, H_2$  respectively. Both hardware are of equal capacity with failure rate  $\lambda_1$  mean repair rate  $w_1$  thus  $2\lambda_1$ , is their combined mean failure rate. Software is similar function which may become inoperable obsolete or need updating by some time with mean failure rate  $\lambda_2$ , so combined failure rate if software is  $2\lambda_2$  and repair / updating rates  $w_2$  each. Separate repairman/ servers are available for hardware and software. If any one of the two hardware fail then it will not be able to run the software hence the system works in reduced capacity or vice-versa i.e. when the software is outdated/ failed then it will not serve any purpose to run it any the hardware, thus the system will also work in reduced state. If both the hardware fails, then the whole of the system is in failed state  $S_2$ , similarly if the software are in failed state then hole of the system is in failed state  $S_5$  if anyone hardware out of any two hardware fails. Then the system works in reduced capacity denoted by state  $S_1$ , similarly if one of the two softwares fails and hardware are good. Then again the system works in reduced state denoted by  $S_4$ , if both the software fail then the system real has the failed state denoted  $S_5$  of either of the hardware and its complementary software fail, then again the system reaches failed state denoted by  $S_3$ , where two separate repairman for hardware and software are available, if the hardware is repaired first then system is transited to state  $S_1$  and if software is repaired first then the system transited to state  $S_4$ .

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**Assumptions and Notations:** The following assumptions and notations are taken: -

1. Separate repair facilities are available for hardware and software.
2. The distributions of failure times and repair times are exponential and general respectively and also different.
3. Failures and repairs are statistically independent.
4. Repair is perfect and repaired system is as good as new one.
5. Nothing can fail when the system is in failed state.
6. The system is discussed for steady-state conditions.
7. Repair facility is immediate.

$\overline{cycle}$  : A circuit formed through un-failed states. m-cycle: A circuit (may be formed through regenerative or non-regenerative / failed state) whose terminals are at the regenerative state m.

$m-\overline{cycle}$  : A circuit (may be formed through un-failed regenerative or non-regenerative state) whose terminals are at the regenerative m state.

$(i \xrightarrow{sr} j)$  : r-th directed simple path from i-state to j-state; r takes positive integral values for different paths from i-state to j-state.

$(\xi \xrightarrow{fff} i)$  : A directed simple failure free path from  $\xi$ -state to i-state.

$V_{m,m}$  : Probability factor of the state m reachable from the terminal state m of the m-cycle.

$V_{m,m}$  : Probability factor of the state m reachable from the terminal state m of the  $m-\overline{cycle}$ .

$R_i(t)$  : Reliability of the system at time t, given that the system entered the un-failed regenerative state 'i' at t = 0.

$A_i(t)$  : Probability of the system in up time at time 't', given that the system entered regenerative state 'i' at t = 0.

$B_i(t)$  : Reliability that the server is busy for doing a particular job at time 't'; given that the system entered regenerative state 'i' at t = 0.

$V_i(t)$  : The expected no. of server visits for doing a job in (0,t] given that the system entered regenerative state 'i' at t = 0.

',' : denote derivative

$W_i(t)$  : Probability that the server is busy doing a particular job at time t without transiting to any other regenerative state 'i' through one or more non-regenerative states, given that the system entered the regenerative state 'i' at t = 0.

$\mu_i$  : Mean sojourn time spent in state i, before visiting any other states;

$$\mu_i = \int_0^\infty R_i(t) dt$$

$\mu_i^1$  : The total un-conditional time spent before transiting to any other regenerative states, given that the system entered regenerative state 'i' at t=0.

$n_i$  : Expected waiting time spent while doing a given job, given that the system entered regenerative state 'i' at t=0;

$$\eta_i = W_i^*(0).$$

$\xi$  : Base state of the system.

$f_j$  : Fuzziness measure of the j-state.

$\lambda_1, \lambda_2$  : Constant failure rate of hosts etc.

$H_i/h_i$  : Unit in full capacity working / failed state of hosts, similarly for servers  $S_1$  &  $S_2$ .

$w_i$  : are constant repair rates of units.

○ Regenerative State, ○ Reduced State, □ Failed State.

Following the above assumptions and notations the Transition Diagram of the system is given in Figure 1.

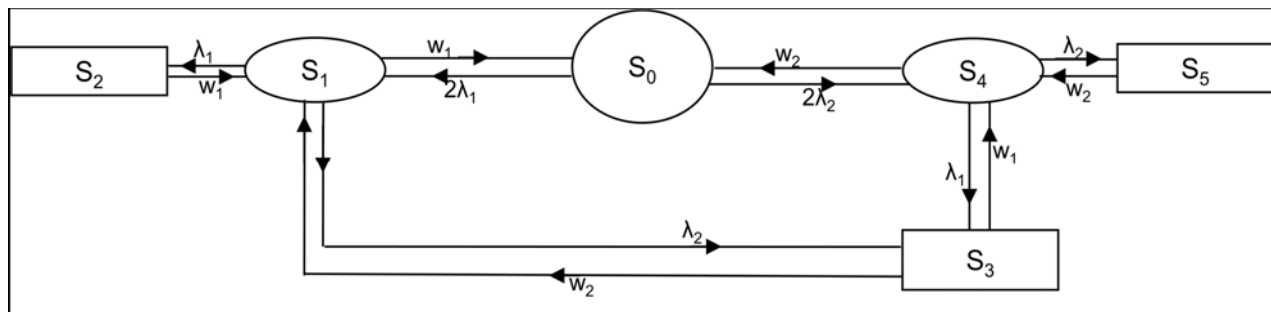


Fig 1

$$S_0 = H_1 H_2 S_1 S_2, \quad S_1 = H_1 h_j S_1 S_2, \quad i \neq j, \quad i, j = 1, 2, \quad S_2 = h_1 h_2 S_1 S_2$$

$$S_3 = h_i H_j S_j S_i, \quad i \neq j, \quad i, j = 1, 2, \quad S_4 = H_1 H_2 S_i S_j, \quad i \neq j, \quad i, j = 1, 2, \quad S_5 = H_1 H_2 S_1 S_2$$

**Table 1:** Simple Paths from vertices

Vertex i	0	1	2	3	4	5
0	(0,1,0) (0,4,0) (0,1,3,4,0)	(0,1) (0,4,3,1)	(0,1,2) (0,4,3,1,2)	(0,1,3) (0,4,3)	(0,4) (0,1,3,4)	(0,4,5) (0,1,3,4,5)
1	(1,0) (1,3,4,0)	(1,2,1) (1,3,4,0,1) (1,0,1) (1,3,1)	(1,2)	(1,3) (1,0,4,3)	(1,0,4) (1,3,4)	(1,0,4,5) (1,3,4,5)
2	(2,1,0) (2,1,3,4,0)	(2,1)	(2,1,2)	(2,1,3) (2,1,0,4,3)	(2,1,0,4) (2,1,3,4)	(2,1,0,4,5) (2,1,3,4,5)
3	(3,1,0) (3,4,0)	(3,1) (3,4,0,1)	(3,1,2) (3,4,0,1,2)	(3,1,3) (3,4,3) (3,1,0,4,3)	(3,4) (3,1,0,4)	(3,0,4,5) (3,4,5)
4	(4,0) (4,3,1,0)	(4,0,1) (4,3,1)	(4,0,1,2) (4,3,1,2)	(4,3) (4,0,1,3)	(4,3,4) (4,5,4) (4,0,4) (4,0,1,3,4)	(4,5)
5	(5,4,0) (5,4,3,1,0)	(5,4,0,1) (5,4,3,1)	(5,4,0,1,2) (5,4,3,1,2)	(5,4,3) (5,4,0,1,3)	(5,4)	(5,4,5)

The possible transitions rates between states along with transition states are shown in Figure 1. Primary, Secondary and Tertiary Circuits associated with the system are given in Table 2.

**Table 2:** Primary, Secondary & Tertiary Circuits at various vertices.

Vertex i	Primary Circuits (CL 1)	Secondary Circuits (CL 2)	Tertiary Circuits (CL 3)	Tetra Circuits (CL 4)
0	(0,1,0) (0,4,0) (0,1,3,4,0)	(1,2,1) (1,3,1) (4,5,4) (4,3,4) (1,2,1) (1,3,1) (3,4,3) (4,5,4)	Nil (3,4,3) Nil (3,1,3) Nil (3,4,3) (4,5,4) Nil	Nil (4,5,4) Nil (1,2,1) Nil (4,5,4) Nil Nil
1	(1,2,1) (1,3,4,0,1)  (1,0,1) (1,3,1)	Nil (3,4,3)  (4,5,4) (0,4,0)  (0,4,0) (3,4,3)	Nil (4,5,4) (4,0,4) Nil (4,5,4) (4,3,4) (4,5,4) (4,3,4) (4,5,4) (4,0,4)	Nil Nil Nil Nil Nil Nil Nil Nil
2	(2,1,2)	(1,0,1) (1,3,1)  (1,0,4,3,1)	(0,4,0)  (3,4,3)  (0,4,0)  (4,3,4) (4,5,4)	(4,5,4) (4,3,4) (4,5,4) (4,0,4) (4,5,4) (4,3,4) Nil Nil
3	(3,1,3) (3,4,3) (3,1,0,4,3)	(1,0,1) (1,2,1) (4,0,4) (4,5,4) (1,0,1) (1,2,1) (0,4,0) (4,5,4)	(0,4,0) Nil (0,1,0) Nil (0,4,0) Nil (4,5,4) Nil	(4,5,4) Nil (1,2,1) Nil (4,5,4) Nil Nil Nil
4	(4,3,4) (4,5,4) (4,0,4)	(3,1,3)  Nil (0,1,0)	(1,0,1) (1,2,1) Nil (1,3,1) (1,2,1)	Nil Nil Nil Nil Nil

	(4,0,1,3,4)	(0,1,0) (1,2,1) (3,1,3)	(1,3,1) (1,2,1) Nil (1,0,1) (1,2,1)	Nil Nil Nil Nil Nil
5	(5,4,5)	(4,3,4) (4,0,4) (4,0,1,3,4)	(3,1,3) (0,1,0)  (0,1,0) (1,2,1) (1,3,1)	(1,0,1) (1,2,1) (1,2,1) (1,3,1) (1,3,1) (1,2,1) Nil Nil

From the table 2, we see that at working state ‘1’ there are maximum number of primary circuits, hence state ‘1’ is the base state.

**Table 3:** Primary, Secondary, Tertiary Circuits w.r.t. the Simple Paths (Base-State ‘1’)

Vertex j	$(1 \xrightarrow{S_r} j): (P0)$	(P1)	(P2)
0	$(1 \xrightarrow{S_1} 0): (1,0)$ $(1 \xrightarrow{S_2} 0): (1,3,4,0)$	(0,4,0) (3,4,3) (4,5,4) (0,4,0)	(4,5,4) (4,3,4) (4,5,4) (4,0,4) Nil (4,5,4) (4,3,4)
1	$(1 \xrightarrow{S_1} 1): (1,2,1)$ $(1 \xrightarrow{S_2} 1): (1,0,1)$ $(1 \xrightarrow{S_3} 1): (1,3,4,0,1)$ $(1 \xrightarrow{S_4} 1): (1,3,1)$	Nil (0,4,0) (0,4,0) (3,4,3) (4,5,4) (4,0,4) (3,4,3)	(4,5,4) (4,3,4) (4,5,4) (4,3,4) (4,5,4) (4,0,4) Nil Nil (4,5,4) (4,0,4)
2	$(1 \xrightarrow{S_1} 2): (1,2)$	Nil	Nil
3	$(1 \xrightarrow{S_1} 3): (1,3)$ $(1 \xrightarrow{S_2} 3): (1,0,4,3)$	(3,4,3) (0,4,0) (4,5,4) (4,3,4)	(4,5,4) (4,0,4) (4,5,4) (4,3,4) Nil Nil
4	$(1 \xrightarrow{S_1} 4): (1,0,4)$ $(1 \xrightarrow{S_2} 4): (1,3,4)$	(0,4,0) (4,5,4) (4,3,4) (3,4,3) (4,5,4) (4,0,4)	(4,5,4) Nil Nil (4,0,4) (4,5,4) Nil (4,5,4) (4,3,4)
5	$(1 \xrightarrow{S_2} 5): (1,0,4,5)$ $(1 \xrightarrow{S_2} 5): (1,3,4,5)$	(0,4,0) (4,5,4) (4,3,4) (3,4,3) (4,5,4) (4,0,4)	(4,5,4) (4,3,4) Nil Nil (4,5,4) (4,0,4) Nil Nil

**Path Probabilities and Mean Sojourn Times.**

$q_{i,j}(t)$  : Probability density function (p.d.f.) of the first passage time from a regenerative state ‘i’ to a regenerative state ‘j’ or to a failed state ‘j’ without visiting any other regenerative state in  $(0,t]$ .

$p_{i,j}$  : Steady state transition probability from a regenerative state ‘i’ to a regenerative state ‘j’ without visiting any other regenerative state.  $p_{i,j} = q_{i,j}^*(0)$  ; where \* denotes Laplace transformation.

**Table 4: Path Probabilities**

$q_{ij}(t)$	$p_{ij} = q_{ij}^*(t)$
$q_{0,1}(t) = 2\lambda_1 e^{-(2\lambda_1+2\lambda_2)t}$	$p_{0,1} = \lambda_1/(\lambda_1+\lambda_2)$
$q_{0,4}(t) = 2\lambda_2 e^{-(2\lambda_1+2\lambda_2)t}$	$p_{0,4} = \lambda_2/(\lambda_1+\lambda_2)$
$q_{1,0}(t) = w_1 e^{-(w_1+\lambda_2+\lambda_1)t}$	$p_{1,0} = w_1/(w_1+\lambda_1+\lambda_2)$
$q_{1,2}(t) = \lambda_1 e^{-(w_1+\lambda_2+\lambda_1)t}$	$p_{1,2} = \lambda_1/(w_1+\lambda_1+\lambda_2)$
$q_{1,3}(t) = \lambda_2 e^{-(w_1+\lambda_2+\lambda_1)t}$	$p_{1,3} = \lambda_2/(w_1+\lambda_1+\lambda_2)$
$q_{2,1} = w_1 e^{-w_1 t}$	$p_{2,1} = 1$
$q_{3,1}(t) = w_2 e^{-(w_1+w_2)t}$	$p_{3,1} = w_2/(w_1+w_2)$
$q_{3,4}(t) = w_1 e^{-(w_1+w_2)t}$	$p_{3,4} = w_1/(w_1+w_2)$
$q_{4,0}(t) = w_2 e^{-(\lambda_1+\lambda_2+w_2)t}$	$p_{4,0} = w_2/(\lambda_1+\lambda_2+w_2)$
$q_{4,3}(t) = \lambda_1 e^{-(\lambda_1+\lambda_2+w_2)t}$	$p_{4,3} = \lambda_1/(\lambda_1+\lambda_2+w_2)$
$q_{4,5}(t) = \lambda_2 e^{-(\lambda_1+\lambda_2+w_2)t}$	$p_{4,5} = \lambda_2/(\lambda_1+\lambda_2+w_2)$
$q_{5,4}(t) = w_2 e^{-w_2 t}$	$p_{5,4} = 1$

It is verified that  $p_{0,1}+p_{0,4} = 1$ ,  $p_{1,0}+p_{1,2}+p_{1,3} = 1$ ,  $p_{2,1} = 1$ ,  $p_{3,1}+p_{3,4} = 1$ ,  $p_{4,3}+p_{4,0}+p_{4,5} = 1$ ,  $p_{5,4} = 1$

**Table 5: Mean Sojourn Times**

$R_i(t)$	$\mu_i = R_i^*(0)$
$R_0(t) = e^{-2(\lambda_2+\lambda_1)t}$	$\mu_0 = 1/2(\lambda_2+\lambda_1)$
$R_1(t) = e^{-(\lambda_1+\lambda_2+w_1)t}$	$\mu_1 = 1/(\lambda_1+\lambda_2+w_1)$
$R_2(t) = e^{-w_1 t}$	$\mu_2 = 1/(w_1)$
$R_3(t) = e^{-(w_1+w_2)t}$	$\mu_3 = 1/(w_1+w_2)$
$R_4(t) = e^{-(\lambda_1+\lambda_2+w_2)t}$	$\mu_4 = 1/(\lambda_1+\lambda_2+w_2)$
$R_5(t) = e^{-w_2 t}$	$\mu_5 = 1/w_2$

**MTSF ( $T_0$ ):** The regenerative un-failed states to which the system can transit (initial state ‘0’), before entering any failed state are: ‘i’ = 0,1,4 taking ‘ $\xi$ ’ = ‘1’.

$$MTSF (T_0) = \left[ \sum_{i,sr} \left\{ \frac{\left\{ \text{pr} \left( \xi \xrightarrow{sr(sff)} i \right) \right\} \mu_i}{\prod_{m_1 \neq \xi} \{1 - V_{m_1 m_1}\}} \right\} \right] \div \left[ 1 - \sum_{sr} \left\{ \frac{\left\{ \text{pr} \left( \xi \xrightarrow{sr(sff)} \xi \right) \right\}}{\prod_{m_2 \neq \xi} \{1 - V_{m_2 m_2}\}} \right\} \right]$$

$$= (V_{0,0}\mu_0 + V_{0,1}\mu_1 + V_{0,4}\mu_4) / (V_{0,0}\mu_0 + V_{0,1}\mu_1 + V_{0,2}\mu_2 + V_{0,3}\mu_3 + V_{0,4}\mu_4 + V_{0,5}\mu_5)$$

**Availability of the System:** The regenerative states at which the system is available are ‘j’ = 0,1, 2 and taking ‘ $\xi$ ’ = ‘1’ the total fraction of time for which the system is available is given by

$$A_0 = \left[ \sum_{j,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow j) \right\} f_j \mu_j}{\prod_{m_1 \neq \xi} \{1 - V_{m_1 m_1}\}} \right\} \right] \div \left[ \sum_{i,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow i) \right\} \mu_i^1}{\prod_{m_2 \neq \xi} \{1 - V_{m_2 m_2}\}} \right\} \right]$$

$$= \left[ \sum_j V_{\xi,j}, f_j, \mu_j \right] \div \left[ \sum_i V_{\xi,i}, f_j, \mu_i^1 \right]$$

$$= (V_{1,1}\mu_1 + V_{1,0}\mu_0 + V_{1,4}\mu_4) / (V_{1,0}\mu_0 + V_{1,1}\mu_1 + V_{1,2}\mu_2 + V_{1,3}\mu_3 + V_{1,4}\mu_4 + V_{1,5}\mu_5)$$

**Busy Period of the Server:** The regenerative states where servers ‘j’ = 1,2,3,4,5 and regenerative states are ‘i’ = 0 to 5, taking  $\xi$  = ‘1’, the total fraction of time for which the servers remains busy is given by

$$B_0 = \left[ \sum_{j,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow j) \right\} n_j}{\prod_{m_1 \neq \xi} \{1 - V_{m_1 m_1}\}} \right\} \right] \div \left[ \sum_{i,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow i) \right\} \mu_i^1}{\prod_{m_2 \neq \xi} \{1 - V_{m_2 m_2}\}} \right\} \right]$$

$$= \left[ \sum_j V_{\xi,j}, n_j \right] \div \left[ \sum_i V_{\xi,i}, \mu_i^1 \right]$$

$$= (V_{1,1}\mu_1 + V_{1,2}\mu_2 + V_{1,3}\mu_3 + V_{1,4}\mu_4 + V_{1,5}\mu_5) / D_1$$

**Expected Number of Inspections by the repair man:** The regenerative states where the repair men join for repair a fresh are 1, 2 to do this job j = 1 the regenerative states are i = 0 to 5, Taking ‘ $\xi$ ’ = ‘1’, the fractional number of visits by the repairmen is given by

$$V_0 = \left[ \sum_{j,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow j) \right\}}{\prod_{k_1 \neq \xi} \{1 - V_{k_1 k_1}\}} \right\} \right] \div \left[ \sum_{i,sr} \left\{ \frac{\left\{ \text{pr}(\xi^{sr} \rightarrow i) \right\} \mu_i^1}{\prod_{k_2 \neq \xi} \{1 - V_{k_2 k_2}\}} \right\} \right]$$

$$= \left[ \sum_j V_{\xi,j} \right] \div \left[ \sum_i V_{\xi,i}, \mu_i^1 \right]$$

$$= (V_{1,1} + V_{1,4}) / D_1$$

**Particular Cases**

$$MTSF (T_0) = (1/D) \left[ (1/4\lambda) + \{ (2w+3\lambda)/4(w+\lambda)(w+2\lambda) \} \right] \left[ \{ 2w/(w+\lambda) \} + \{ \lambda(w+4\lambda)/(w+4\lambda) \} \right]$$

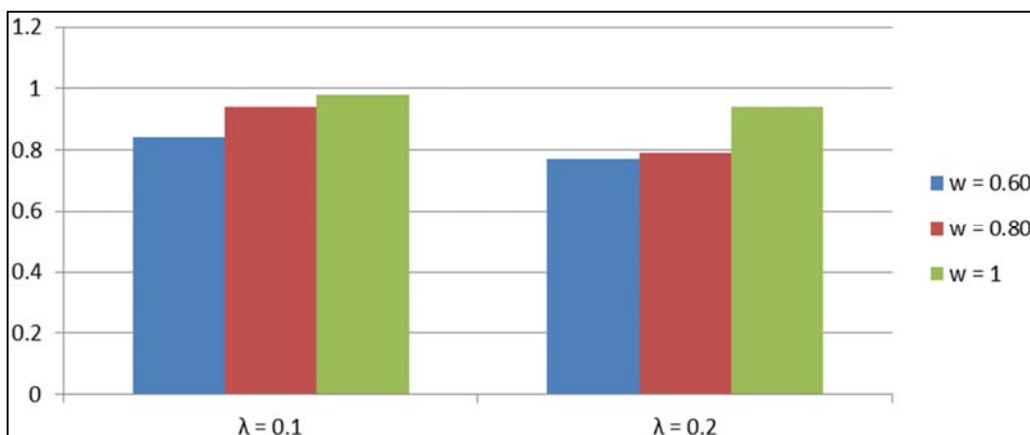
$$= (w+\lambda) - \lambda + \{ \lambda w(2w+3\lambda)/(2w+\lambda)^2(w+\lambda) \}$$

Where  $D = V_{0,0}\mu_0 + V_{0,1}\mu_1 + V_{0,2}\mu_2 + V_{0,3}\mu_3 + V_{0,4}\mu_4 + V_{0,5}\mu_5$

**Table 6:** MTSF Table

T <sub>0</sub>	w = 0.60	w = 0.80	w = 1
λ = 0.1	0.84	0.94	0.98
λ = 0.2	0.77	0.79	0.94

**MTSF Graph**



**Fig 2**

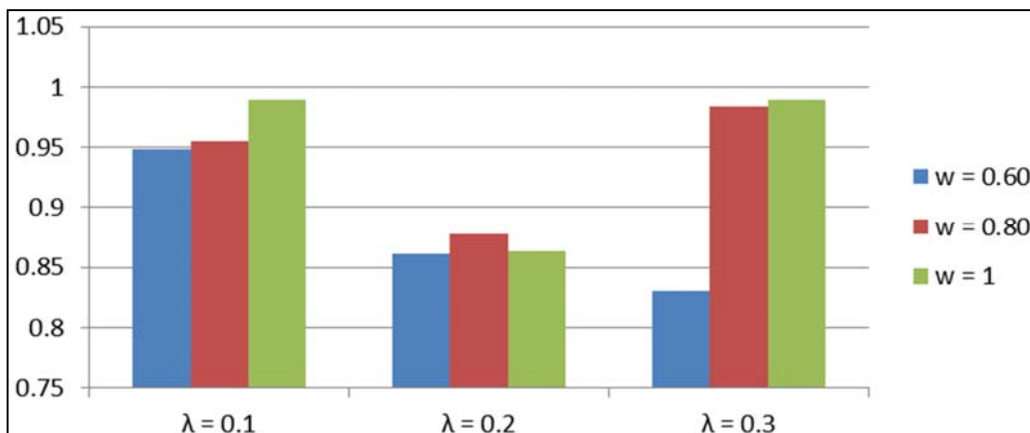
**Availability of the System (A<sub>0</sub>) =** [(1/D<sub>1</sub>)(V<sub>1,1</sub>μ<sub>1</sub>+V<sub>1,0</sub>μ<sub>0</sub>+V<sub>1,4</sub>μ<sub>4</sub>)]

Where D<sub>1</sub> = (V<sub>1,0</sub>μ<sub>0</sub>+V<sub>1,1</sub>μ<sub>1</sub>+V<sub>1,2</sub>μ<sub>2</sub>+V<sub>1,3</sub>μ<sub>3</sub>+V<sub>1,4</sub>μ<sub>4</sub>+V<sub>1,5</sub>μ<sub>5</sub>)  
 = (1/D<sub>1</sub>)[{1/(w+2λ)}+{1(w+λ)(2w+3λ)/4λ(w+λ)+(2w+3λ)-w(w+2λ)}]  
 [{w/(w+2λ)}+{λw(w+4λ)/2(w+2λ)(w+2λ)(w+4λ)-λ(w+4λ)-λ(w)}]  
 +{4w/(2w+3λ)(w+2λ)}+{λ/2(w+2λ)}][(w+4λ)/(w+λ)(w+4λ)-λ(w+2λ)]  
 {1/(w+λ)}  
 = (1/D<sub>1</sub>)[{1/(w+2λ)}+{1+(w+λ)(2w+3λ)/4λ(w+λ)+(2w+3λ)-w(w+2λ)}]  
 [w+{λw(w+4λ)/2(w+2λ)(w+4λ)-2λ(w+3λ)} +{4w/(2w+3λ)}+{λ(w+4λ)  
 (w+λ)(2w+3λ)-w(w+2λ)/2(w+λ)(w+4λ)-λ(w+2λ)(w+λ)}]  
 = [1/(w+2λ)1+{1+(w+λ)(2w+3λ)/4λ(w+λ)+(2w+3λ)-w(w+2λ)}]  
 [w+{λw(w+4λ)/2(w+2λ)(w+4λ)-2λ(w+3λ)}+{4w/(2w+3λ)}+{λ(w+4λ)  
 (w+λ)(2w+3λ)-w(w+2λ)/2(w+λ)(w+4λ)-λ(w+2λ)(w+λ)(2w+3λ)}+λ  
 +{λ(w+λ)(w+4λ)/(w+λ)(w+4λ)-λ(w+2λ)}+{wλ(w+λ)/4(w+2λ)<sup>3</sup>}  
 -{w<sup>2</sup>λ/4(w+2λ)<sup>2</sup>}+{λw/(w+λ)(2w+3λ)-w(w+2λ)}+{λ<sup>2</sup>/(w+λ)(w+4λ)-λ(w+2λ)}]

**Table 7:** Availability Table

A <sub>0</sub>	w = 0.60	w = 0.80	w = 1
λ = 0.1	0.948	0.955	0.989
λ = 0.2	0.861	0.878	0.863
λ = 0.3	0.830	0.984	0.989

**Availability of the System Graph**



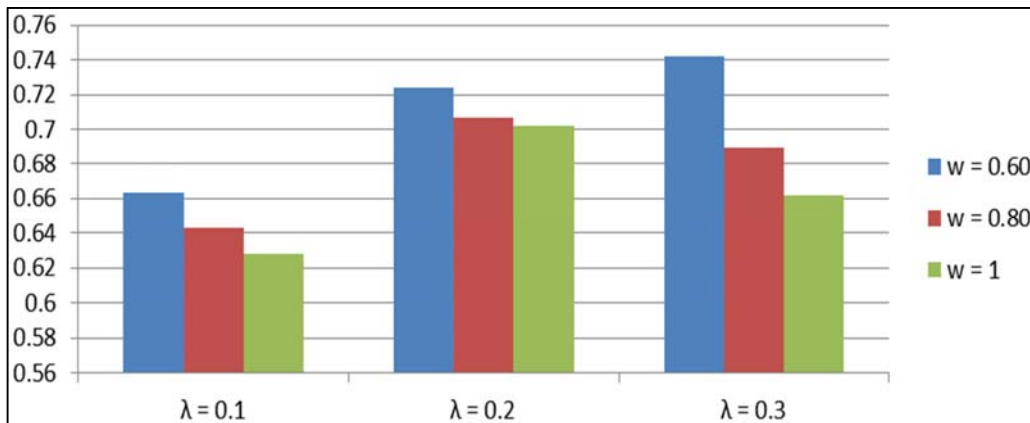
**Fig 3**

**Busy Period of the Server (B<sub>0</sub>)** =  $[1-(V_{1,0}\mu_0/D_1)]$ , =  $[1-(\mu_0/D_1)]$ , =  $[1-(1/4\lambda D_1)]$

**Table 8:** Busy Period of the Server Table

B <sub>0</sub>	w = 0.60	w = 0.80	w = 1
λ = 0.1	0.663	0.643	0.628
λ = 0.2	0.7241	0.707	0.702
λ = 0.3	0.742	0.689	0.662

**Busy Period of the Server Graph**



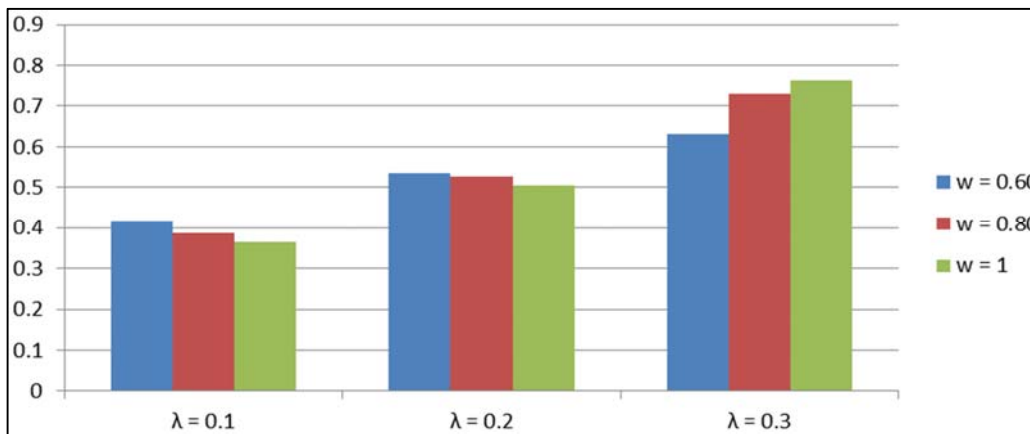
**Fig 4**

**Expected Number of Server's Visits (V<sub>0</sub>)** =  $(1/D_1) [1 + \{4w/(2w+3\lambda)(w+2\lambda)\} + \{\lambda/2(w+2\lambda)\} \{(w+4\lambda)/(w+\lambda)(w+4\lambda)-\lambda(w+2\lambda)\} \{1/(w+\lambda)(2w+3\lambda)\}]$

**Table 9:** Expected Number of Server's Visits Table

V <sub>0</sub>	w = 0.60	w = 0.80	w = 1
λ = 0.1	0.417	0.388	0.366
λ = 0.2	0.536	0.528	0.504
λ = 0.3	0.631	0.730	0.763

**Expected Number of Server Visits Graph**



**Fig 5**

**Profit Function**

= A<sub>0</sub>R<sub>0</sub> - K<sub>0</sub>B<sub>1</sub> - K<sub>1</sub> (No. of Visits)

= 500 A<sub>0</sub> - 200 B<sub>0</sub> - 100 (No. of Visits)

Taking R<sub>0</sub> = 500 K<sub>0</sub> = 200 K<sub>1</sub> = 100

**Table 10:** Profit Function Table

	w = 0.60	w = 0.80	w = 1
λ = 0.1	299.70	310.10	332.30
λ = 0.2	232.08	244.80	240.70
λ = 0.3	203.50	281.20	310.80

### Profit Function Graph

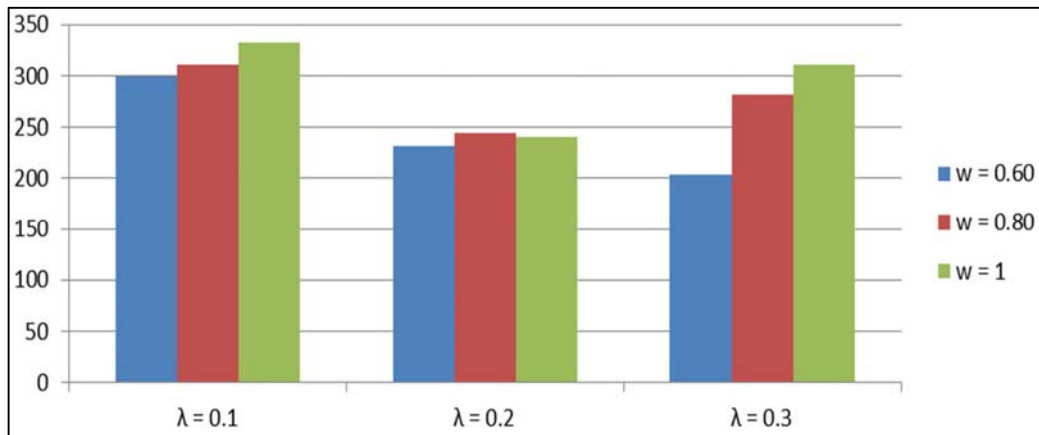


Fig 6

**Conclusion:** From above tables and graph we see that the result obtained using Regenerative Point Graphical Technique is practical, trendy and same as obtained by using Regenerative Point Technique and other techniques. But in Regenerative Point Graphical Technique, we obtained the results very easily and quickly without writing any state equations and without any cumbersome procedures, long calculations and simplifications. Regenerative Point Graphical Technique is applied to study the behavior and profit analysis of various process industries like soap, soft drink and dairy plant, paper industry, soap industry etc. It is hoped that the Regenerative Point Graphical Technique for the analysis of the system will be very helpful to the managements, manufactures and the personal engaged in reliability engineering and working for the behavior and profit analysis of stochastic systems.

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